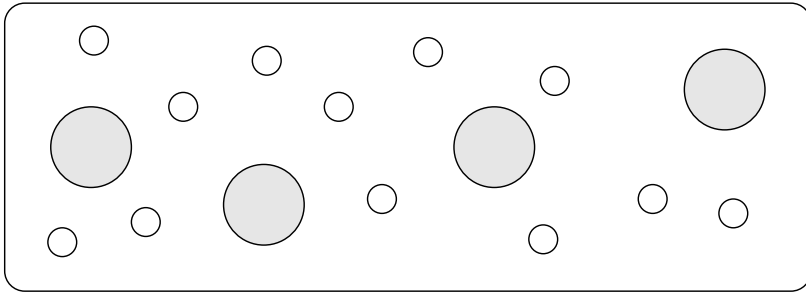


# Error repair



Good programming languages are designed with a relatively large “distance” between syntactically correct programs, to increase the likelihood that conceptual mistakes are caught as syntactic errors.

Error repair usually occurs at two levels:

**Local:** repairs mistakes with little global import, such as missing semicolons and undeclared variables.

**Scope:** repairs the program text so that scopes are correct. Errors of this kind include unbalanced parentheses and begin/end blocks.

Repair actions can be divided into *insertions* and *deletions*. Typically the compiler will use some lookahead and backtracking in attempting to make progress in the parse. There is great variation among compilers, though some languages (PL/C) carry a tradition of good error repair. Goals of error repair include:

1. No input should cause the compiler to collapse.
2. Illegal constructs are flagged.
3. Frequently occurring errors are repaired gracefully.
4. Minimal stuttering or cascading of errors.

LL-style parsing lends itself well to error repair, since the compiler uses the grammar’s rules to *predict* what should occur next in the input.

# Augmenting recursive descent parsers for error recovery

Recursive and LL parsers are often called *predictive*, because they operate by predicting the next step in a derivation.

Suppose the parser is operating in procedure  $A$  for some nonterminal  $A$ . If an error occurs, it seems reasonable to recover by skipping to a symbol that could follow  $A$ , and then return.

$$\begin{array}{l}
 \mathbf{E} \rightarrow \mathbf{T} E' \\
 \hline
 E' \rightarrow + \mathbf{T} E' \\
 E' \rightarrow - \mathbf{T} E' \\
 \quad | \quad \lambda \\
 \hline
 \mathbf{T} \rightarrow \mathbf{F} T' \\
 T' \rightarrow * \mathbf{F} T' \\
 T' \rightarrow / \mathbf{F} T' \\
 \quad | \quad \lambda \\
 \mathbf{F} \rightarrow (\mathbf{E}) \\
 \quad | \quad \mathbf{a}
 \end{array}$$

	First	Follow
<b>E</b>	{ (, a }	{ ), \$ }
$E'$	{ +, - }	{ ), \$ }
<b>T</b>	{ (, a }	{ +, -, ), \$ }
$T'$	{ *, / }	{ +, -, ), \$ }
<b>F</b>	{ (, a }	{ *, /, +, -, ), \$ }

**Procedure**  $E'(StopSet)$

```

if ( $LookAhead(+)$ ) then
    call  $Expect(+)$ 
    call  $T(\{+, -\} \cup StopSet)$ 
    call  $E'(StopSet)$ 
else
    if ( $LookAhead(('$, ')')$ ) then
        return ()
    else
        call  $ErrorRecover(StopSet)$ 
    fi
fi
end

```